

Amendments to the Specification:

Please replace the paragraph beginning at page 4, line 15 with the following amended paragraph:

FIG. 1 shows the system architecture for performing the scheduling method of the present invention, wherein the number of flows in the system is assumed to be N. The packet from flow i is processed by the present scheduling method and output through the output queue 11. The output queue 11 has a plurality of windows 111, each window 111 having a size of W. The flows 1, 2, ...N have weights of $W_1^*, W_2^*, \dots, W_N^*$, respectively. Based on the size W of each window 111, the weight W_i^* of each flow i can be normalized to be $w_i = W \times W_i^* / (W_1^* + W_2^* + \dots + W_N^*)$.

Please replace the paragraph beginning at page 5, line 8 with the following amended paragraph:

If step S11 determines that the credit c_i of flow i is larger than the size of packet P_i , the packet P_i is placed into one of the windows 111, for example the kth window (window k of FIG. 3) [[k]], pointed by the window index d_i for being output through the output queue 11 (step S13). In step S14, the size of the packet P_i is subtracted from the credit c_i , and the window index w_i and the credit c_i are updated; i.e., the updated window index w_i and the credit c_i are written into the table 12 (step S14).

Please replace the paragraph beginning at page 5, line 15 with the following amended paragraph:

The packets placed in the windows 111 of the output queue 11 are[[is]] sequentially output. As shown by the packet departure process in FIG. 1, when all packets have been pushed out and the window 111 is empty, the table 12 is[[11]] updated to have its initial values.